

Next Presentation begins at 10:10

#### Advanced Security Evaluation of Network Protocols

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providing security.





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# Advanced Security Evaluation of Network Protocols

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- Today I'll talk about evaluation of (proprietary) network protocols
  - Why is it necessary ?
  - What is the typical methodology ?
  - What can be done better ?









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# The story of Heartbleed



Heartbleed

• aka. CVE-2014-0160



# The story of Heartbleed

- Infosecurity INTELLIGENT
- Heartbeats needed for DTLS (TLS over UDP) to keep NAT states active.
- Heartbeats are also present in TLS (over TCP), even thou they are unnecessary.
- Heartbeats doesn't need to have a payload, but they have.
- The payload doesn't need to be variable in length, but it is.



# The story of Heartbleed

- If heartbeats include a payload of variable length, at least the length should be checked.
- But the length isn't checked, resulting in the ability to read from the following memory segment.
- Wouldn't be that much of a problem, if OpenSSL would use standard memory management instead of its own.

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# The story of SNMPv3 HMAC



• SNMPv3 HMAC Bug

• aka. CVE-2008-0960



# The story of SNMPv3 HMAC

- SNMPv3 supports HMAC authentication.
- The HMAC can be of variable (user defined) length.
- Even a length of one byte could be chosen.
- Resulting in an authenticator with 256 possible values
   => Easy to brute force

byte (8-bits



# The story of SNMPv3 HMAC



• HMACs of dynamic length might be a good idea, but please define a minimal (secure) length!





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# The story of Ping of Death

- Ping of Death
- Originally appeared in 1996
- Buffer overflow with ICMP packet bigger than 2<sup>16</sup> bytes.
- Results in Denial of Service.
- Effected large amount of Operating Systems, including Unix, Linux, Mac and Windows.



# The story of Ping of Death

- Ping of Death reappeared in 2013 on Windows systems.
- This time in ICMPv6.
- aka. MS13-065
- Exact same vulnerability, 17 years later.









# The story of the CTL

- Cisco VoIP Certificate Trust List
- Not a protocol per se, but proprietary file format used in combination with proprietary network protocol.



# The story of the CTL

- Certificate Trust List is fetched during provisioning of VoIP Phones and stored as root of trust.
- The initial CTL is blindly trusted (you have to trust your root of trust, don't you?).
- Nobody noticed the Cisco IP Communicator (the VoIP softphone) deleted the CTL on every shutdown.
- => see "All Your Calls Are Still Belong to Us: How We Compromised the Cisco VoIP Crypto Ecosystem" for details.









# Vulnerability in ASN.1 libraries

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- A lot of them have appeared in the past.
- To mention a few:
  - CVE-2003-0543
  - CVE-2003-0544
  - CVE-2003-054
  - MS04-007
  - CVE-2005-1730
  - CVE-2005-1935



# Vulnerability in ASN.1 libraries

- They affect all tools using the library to parse ASN.1.
- Some of them allow remote code execution.
- Hard to spot, as ASN.1 is complex and libraries should be well tested.
- Ironically ASN.1 libs are used to keep you save from this kind of bugs.





# Vulnerability in ASN.1 libraries

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- Don't blindly rely on protocol parsing libraries.
- Even if your service is using ASN.1, testing on the protocol level still is needed.







# All beginnings are difficult

- How would you start analyze any protocol ?
  - Right, RTFM.
- How would you start analyze an undocumented protocol ?
  - Not so easy.



#### An Example

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#### • Lets exercise this on an example.

- I've chosen an undocumented, proprietary protocol that has crossed my path in the past.
- Was used by a Fat Client and a Java applet.
- First we'll ask our old friend wireshark for help.





eq=0 Win=1460 (CK] Seq=0 Acl (CK] Seq=1 Acl (CK] Seq=1 Acl (CK] Seq=1 Acl (CK] Seq=1 Acl (CK] Seq=2 Acl
CK] Seq=0 Acl Seq=1 Ack=1 W: CK] Seq=1 Ack Seq=1 Ack=137 ACK] Seq=1 Ack Seq=137 Ack=2 ACK] Seq=2 Ack
Geq=1 Ack=1 W: CK] Seq=1 Ack Geq=1 Ack=137 ACK] Seq=1 Ack Geq=137 Ack=2 ACK] Seq=2 Ack
CK] Seq=1 Ack Seq=1 Ack=137 CK] Seq=1 Ack Seq=137 Ack=2
Geq=1 Ack=137 ACK] Seq=1 Ack Geq=137 Ack=2 ACK1 Seq=2 Ack
CK] Seq=1 Ack Seq=137 Ack=2 ACK1 Seq=2 Ack
Seq=137 Ack=2
CK1 Sen=2 Ack
ACIC] Deq-2 ACI
eq=137 Ack=30
\CK] Seq=137 /
eq=305 Ack=10
\CK] Seq=161 /
eq=305 Ack=3;
\CK] Seq=329 /
eq=305 Ack=30
\CK] Seq=305 /
eq=361 Ack=30
\CK] Seq=306 /
\CK] Seq=361 /
eq=417 Ack=38
\CK] Seq=385 /
eq=417 Ack=4

#### The data

- The protocol is TCP based.
- It uses port 8401.
- Lets take a look at the transferred data.

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0-

#### The data



- Some weird binary stuff in the beginning.
  - Includes user authentication.
  - That's going to be the fun part (-;
- XML payload later on.







• Authentication always is interesting.



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Stream Content-



### Protocol fields

- Does look like a packet header, followed by some payload.
- How do I know?
  - Typical 4 byte integer values (Big Endian, aka. Network byte order) at the beginning.
  - ASCII payload in the end.



#### Stream Content-

0000000	00 0	0 00 0	00 00	00 00	) 80	00	00	00	00	00	00	00	00				
00000010	00 0	0 04 3	32 00	00 00	01 (	00	00	09	5f	00	00	00	68	2			ŀ
00000020	6b 2	e 4e 2	26 46	43 62	2 50	59	53	76	24	48	42	71	21	k.N&	=CbP	YSv\$HE	3q !
00000030	25-3	0 3d 7	79 2c	72 46	9 5f	50	63	45	40	5a	57	73	57	%0=y	, rN	PcE@ZV	NsV
00000040	4f 7	0 42 5	5c 31	55 70	) 5d	49	Зf	7d	30	23	56	48	47	OpB\.	1Up]	I?}0#\	VH(
00000050	3a 3	c 54 2	25 36	48 51	. 30	53	35	58	47	45	7c	37	36	. <t%∣< td=""><td>6HQ0</td><td>S5XGE</td><td>76</td></t%∣<>	6HQ0	S5XGE	76
00000060	78 6	0 6e 4	18 7c	48 79	) 5c	61	72	2f	5b	7f	46	4a	29	x`nH	Hy\	ar/[.F	EJ)
00000070	3a 6	9 48 5	5a 44	33 31	. 27	43	23	7a	21	54	31	6a	25	:iHZI	D31 '	C#z!TI	۱j۹
00000080	70 4	4 3b 3	31 00	00 00	00 (									pD;1			
00000000 00																	
00000001 00	00 2e	00 00	00 (	0c 00	- 00	00	00 0	01 (	cd (	e5 .	72 (	90				<b>r</b> .	
00000011 00	04 2f	00 00	00 (	2f 00	00	00	80 (	00 (	90 (	00 (	90 (	91	/ .	/.			
00000021 cd	e5 72	00 00	04:	2f 00	00	00	01 (	00 (	00 (	09 (	50 (	90	r.	/.		โ.	
00000031 00	00 68	39 50	: 50 !	50 75	72	5c	5d (	69 0	37 🔅	39 (	56 2	2a	h9	)∖PPu	[//r	i79f*	
00000041 54	64 Ge	5d 65	5 56 0	32 6b	2f	32	2c (	64 🕻	26	78 :	3b 🕻	35 1	[dn	eV2k	/2,0	l&x5	
00000051 45	3e 32	20 30	79 :	3c 2a	64	6c	4a 🕻	27 0	72 🔅	33 4	4b (	Sb I	E>2	=y<*	dlJ	' n3Kk	
00000061 4f	24 Sc	62 51	. 52 (	61 6a	7b	2e	76 4	4c (	6b -	40 4	47 (	69 I	D\$∖b	QRaj	{.vl	_k@Gi	
00000071 3b	36 36	76 75	5 6f !	5d 28	7b	65	44 🖯	73 0	75 🔅	31 4	42 0	78	; 66\	/uo] (	{eDs	sulBx	
00000081 64	4b 60	62 28	3 4e !	53 54	63	7c	58 0	3e J	78	62 (	59 4	4c (	dK`b	(NST	c X:	>xbi∟	
00000091 49	65 22	54 35	5 2b -	45 00	00	00	00 (	00 (	90	00 0	30 (	90 3	Ie"1	<u>5+</u> E.		0.	







- How to identify those fields?
- Lets first check for the obvious ones:



# Protocol fields

<sub>F</sub> Stream Conter	nt—																		
00000000	00	00	00	09	00	00	00	80	00	00	00	00	00	00	00	00			
00000010	00	00	04	32	00	00	00	01	00	00	09	5f	00	00	00	68	2	h	
00000020	6b	2e	4e	26	46	43	62	50	59	53	76	24	48	42	71	21	k.N&FCbP	YSv\$HBq!	
00000030	25	30	Зd	79	2c	72	4e	5f	50	63	45	40	5a	57	73	57	%0=y,rN_	PcE@ZWsW	
00000040	4f	70	42	5c	31	55	70	5d	49	Зf	7d	30	23	56	48	47	OpB\1Up]	I?}0#VHG	
00000050	За	Зc	54	25	36	48	51	30	53	35	58	47	45	7c	37	36	: <t%6hq0< td=""><td>S5XGE 76</td><td></td></t%6hq0<>	S5XGE 76	
00000060	78	60	6e	48	7c	48	79	5c	61	72	2f	5b	7f	46	4a	29	x`nH Hy\	ar/[.FJ)	
00000070	Зa	69	48	5a	44	33	31	27	43	23	7a	21	54	31	6a	25	:iHZD31'	C#z!Tlj%	
00000080	70	44	Зb	31	00	00	00	00									pD;1		

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#### Protocol header

- The other fields are not so obvious.
- When looking at a series of packets and the associated answers, other fields such as Type, Command and Sequence No. can be identified.
- Finally we can guess the packet header:





#### **Protocol header**







 By triggering different actions in the client and carefully observing the produced traffic, Type and Subtype values can be identified.





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- \$ openssl asn1parse -inform der -in test.bin
  - 0:d=0 hl=2 l= 92 cons: SEQUENCE
  - 2:d=1 hl=2 l= 13 cons: SEQUENCE
  - 4:d=2 hl=2 l= 9 prim: OBJECT :rsaEncryption
  - 15:d=2 hl=2 l= 0 prim: NULL
  - 17:d=1 hl=2 l= 75 prim: BIT STRING
  - 94:d=0 hl=2 l= 0 prim: EOC



![](_page_40_Picture_9.jpeg)

\$ openssl rsa -inform der -in test.bin -pubin -text
Public-Key: (512 bit)

Modulus:

00:c8:4b:c9:ee:7f:de:99:ac:5d:d0:c6:a1:cc:1c:

40:e7:f7:6c:44:50:7d:09:81:a5:71:76:0c:9a:97:

Od:ee:56:a2:fc:74:ce:d1:f3:68:ae:16:c2:a2:23:

6f:06:c6:b2:0d:70:bb:99:fc:45:79:8b:d2:5b:a7:

d6:49:9a:d2:29

Exponent: 65537 (0x10001)

![](_page_41_Picture_9.jpeg)

![](_page_41_Picture_10.jpeg)

# Interesting values

![](_page_42_Picture_1.jpeg)

Туре	Subtype	Content
0x9	0x9F6	Server PubKey
0x2F	0x9FB	Session Key
0x9	0xA00	Login Data

![](_page_42_Picture_3.jpeg)

# What's happening here ?

- The server transmits its public key.
- Client uses the public key to encrypt the session key.
- Session key is used to encrypt login data.

![](_page_43_Picture_4.jpeg)

![](_page_43_Picture_5.jpeg)

#### But ....

![](_page_44_Picture_1.jpeg)

• The public key sent by the server is never validated.

• How should it be, its no certificate, right?

• .... goto fail;

![](_page_44_Picture_5.jpeg)

![](_page_44_Picture_6.jpeg)

![](_page_45_Picture_1.jpeg)

![](_page_45_Picture_2.jpeg)

\$ python mitm.py
Got client, opening outgoing socket
outgoing socket established

• • • • • •

Found Pub-Key of len 512 Generating new Pub-Key

![](_page_46_Picture_7.jpeg)

![](_page_46_Picture_8.jpeg)

#### Found Session Key 'f8ab5431b0cd73a7'

• • • • • • • • •

#### Found Username 'ernw test'

\*\*\*\*\*

Found Random 'e3b9fc671be3a307'

![](_page_47_Picture_11.jpeg)

![](_page_47_Picture_12.jpeg)

![](_page_47_Picture_13.jpeg)

#### Found Encoded Password

#### '38382ac3b3b2e04ff0513560801af46e9c05e3f8'

![](_page_48_Picture_6.jpeg)

![](_page_48_Picture_7.jpeg)

# The last hurdle

• The transmitted password is not encrypted, but encoded...

```
def decode(data):
    out = []
    data = [ ord(i) for i in data.decode("cp1252") ]
    for i in xrange(len(data)):
        out += [32 + seed[i % 16].index(data[i])]
        return "".join( \
            [chr(i).encode("cp1252") for i in out] )
```

![](_page_49_Picture_3.jpeg)

![](_page_50_Picture_0.jpeg)

![](_page_50_Picture_1.jpeg)

![](_page_50_Picture_2.jpeg)

![](_page_51_Picture_0.jpeg)

![](_page_51_Picture_1.jpeg)

![](_page_52_Picture_0.jpeg)

- Secure protocol design is hard.
- Secure protocol implementation is even harder.
- To avoid security issues with the design as well as the implementation one should always review them from an attackers point of view:

![](_page_52_Picture_4.jpeg)

# Review the design

- Is the protocol authenticated?
- If so, is the authentication data encrypted, not just encoded (think of previous example or ROT13)
- If asymmetric crypto is used, are the public keys validated?

![](_page_53_Picture_4.jpeg)

![](_page_53_Picture_5.jpeg)

# Review the implementation

- Are common pitfalls on the programming language level avoided? (Integers overflows)
- Are flaws on the data representation level avoided?
  - Length fields checked for the actual amount of data present
  - Length fields checked for the buffer size available
- Are logical flows in the protocol validated for the users authorization and correct order?

![](_page_54_Picture_6.jpeg)

![](_page_54_Picture_7.jpeg)

### Some last words

- Even today a lot of bad and worse protocols are in use.
- So please, do evaluate more protocols.
- Even more so, if they are used by a huge amount of software (think of SSL).
- Don't be afraid of proprietary protocols, most of the time there is a reason for them not being documented.

![](_page_55_Picture_5.jpeg)

![](_page_55_Picture_6.jpeg)

# There's never enough time...

![](_page_56_Picture_1.jpeg)

![](_page_56_Picture_2.jpeg)

![](_page_56_Picture_3.jpeg)